

DIFFERENTIAL CALCULUS OF EQUATIONALLY DEFINED FUNCTIONS BY WAY OF POLYNOMIAL EXPANSIONS

Francesca Schremmer, *West Chester University* (700 S. High St., West Chester, PA 19383)
Alain Schremmer, *Community College of Philadelphia* (1700 Spring Garden St., Philadelphia, PA 19130).

Supported in part by NSF Grant USE-8814000

0. Introduction. A calculus for "*just plain folks*", «Workshop, 1986 #35», need not be based on the notion of limits. Indeed, and it is an idea going back at least to «Lagrange, 1797 #17», functions can be studied locally from their jets, that is from their best polynomial approximations, *obtained a priori*. In «Schremmer, In Press #23», we sketched such an approach in the case of polynomial functions and in «Schremmer, In preparation #24», we will discuss rational functions and show how, in this particular case, we can even derive a certain amount of *global* information from a small number of *local* investigations.

Here, after briefly recapitulating the main features of Lagrange's approach, we discuss how it applies to functions defined by functional equations, algebraic and differential.

1. Lagrange's approach. To **expand** a function f near a given point x_0 , we first **localize** it, that is we get a form in which the terms are in descending order of magnitude, by expressing $f(x_0 + h)$, the value of the function f near x_0 , as a polynomial function $F_{x_0}(h)$ plus a remainder $R_{x_0}(h)$

$$f(x_0 + h) = F_{x_0}(h) + R_{x_0}(h),$$

where $F_{x_0}(h) = A_0 + A_1h + A_2h^2 + A_3h^3 + \dots + A_nh^n$ and $R_n(h) = h^n \cdot o[1]$ so that $P_{x_0}(h)$ is the **best polynomial approximation of degree n** of $f(x_0 + h)$. For all practical purposes, we shall just write $f(x_0 + h) = F_{x_0}(h) + \dots$

We take the differential calculus to consist of "*the techniques used to find out certain properties of functions*" «Gleason, 1967 #34». The degree n of the approximation that we require depends on the nature of the required information. For instance, from a *qualitative* viewpoint, the **sign** of f near x_0 is given by the least non-zero approximation (usually the best constant approximation), the **variance** is given by the least non-constant approximation (usually the best affine approximation) and the **concavity** is given by the least non-affine approximation (usually the best quadratic approximation).

From a *quantitative* viewpoint, the i^{th} **derivative** of f is *defined* as the function $f^{(n)}$ whose value at x_0 is $i! A_i$ which gives:

$$f(x_0 + h) = f(x_0) + f'(x_0)h + \frac{f''(x_0)h^2}{2} + \frac{f^{(3)}(x_0)h^3}{3} + \dots + \frac{f^{(n)}(x_0)h^n}{n!} + \dots$$

For instance, $[x^n]' = (n-1)x^{n-1}$ because $(n-1)x_0^{n-1}$ is the coefficient of h in the (binomial) expansion of $(x_0 + h)^n$. Similarly, to obtain the derivative of $[f * g]$, where $*$ is any operation, we take the coefficient of h in the expansion of $[f * g](x_0+h)$.

That we recover the **Taylor expansion** of f should lead us to expect that C^n functions are amenable to Lagrange's approach and, in fact, the statement that "*all decent functions have continuous derivatives*" translates into "*all decent functions are practically (polynomial)*" «Gleason, 1967 #34».

The first question then is how to find the polynomial approximation. By analogy with arithmetic, it is, in the case of polynomial functions, by truncation of the high powers near 0 and of the low powers near ∞ and, in the case of rational functions, by division in ascending powers near 0 and in descending powers near ∞ . Anywhere inbetween requires that we first set $x = x_0 + h$.

2. Functional equations. A distinction worth making at the outset, but in fact usually not emphasized, is that, when we write $f(x) = -3x^3 - 2x + 4$, we are defining the function f by the finite algorithm that gives the value $f(x)$ at any point x while, when we write $g(x) = \sqrt{x}$, we are defining the function g as solution of a functional equation, $g^2(x) = x$, without giving any algorithm for solving this equation and computing the value $g(x)$. It is only in the first case that a function can truly be equated with a machine. Another fact not usually stressed is that, in most cases and even in that of rational functions, the algorithm only gives *approximate* values¹.

A practical consequence is that, when given a function such as $f(x) = \sqrt[3]{\frac{x-1}{\sqrt{x^2+1}}}$, the students rarely realize that the first thing to do is to find the functional equation of which it is the solution that is, with due regard to sign considerations,

$$f^6(x) = \frac{(x-1)^2}{x^2+1}.$$

At this point however, we can find the value of the derivative at 0 faster and more reliably than by evaluating the derivative obtained by the usual rules. To get the Best Affine Approximation near 0, we set $f_0(x) = A_0 + A_1x + \dots$ and substitute in the functional equation:

$$\begin{aligned} [A_0 + A_1x + \dots]^6 &= \frac{(x-1)^2}{x^2+1} \\ A_0^6 + 6 A_1^5x + \dots &= \frac{1 - 2x + \dots}{1 + \dots} \\ &= 1 - 2x + \dots \end{aligned}$$

¹ Even in arithmetic, students are rarely given the opportunity to realize, for instance, that 5^2+3 is a constructive template even if the algorithm is not actually given but that $\sqrt{5}$, or for that matter $\frac{5}{3}$ or even $5-3$, is nothing but the name given a priori to the solution of an equation, $x^2 = 5$, (resp. $3x = 5$ or $x + 3 = 5$), *should it exist*, and that, in fact, we can usually only approximate the solution. Thus, the distinction between arithmetic and algebra is somewhat counterproductive if not disingenuous.

Identifying the coefficients gives $A_0 = 1$ and $A_1 = -\sqrt[5]{\frac{1}{3}}$ so that $f_0(x) = 1 - \sqrt[5]{\frac{1}{3}}x + \dots$ and $f'(0) = -\sqrt[5]{\frac{1}{3}}$. The equation of the tangent at the origin is, of course, $t_0(x) = -\sqrt[5]{\frac{1}{3}}x + 1$.

To obtain the derivative of $f(x)$, we need the coefficient of h at x_0 . We localize:

$$f(x_0 + h)^6 = \frac{(x_0 + h - 1)^2}{(x_0 + h)^2 + 1}$$

and expand

$$(A_0 + A_1h + \dots)^6 = \frac{(x_0-1)^2 + 2(x_0-1)h + \dots}{x_0^2 + 1 + 2x_0h + \dots}$$

Dividing in ascending powers, we get

$$A_0^6 + 6A_1A_0^5h + \dots = \frac{(x_0 - 1)^2}{x_0^2 + 1} + \frac{3x_0^2 - 2x_0 - 1}{x_0(x_0^2 + 1)}h + \dots$$

$$\text{Identifying the coefficients gives } A_0 = \sqrt[6]{\frac{(x_0 - 1)^2}{x_0^2 + 1}} \text{ and } A_1 = \frac{3x_0^2 - 2x_0 - 1}{6x_0(x_0^2 + 1)}$$

Observe that getting the second derivative would not be that much more difficult.

12. EXPONENTIAL AND LOGARITHM FUNCTIONS. There are at least three approaches to the exponential function a^x . The approach most currently favored these days is to begin by introducing the notion of integral in the middle of the differential calculus for the sole purpose of defining e^x as $\left[\int \frac{dx}{x} \right]^{-1}$, the inverse of the indefinite integral of the reciprocal function!

A more natural way would be to extend the notion of power to irrational exponents and introduce a^x as limit of a^{s_n} , where s_n is rational and approaches x as n approaches ∞ . Unfortunately, establishing the usual computational rules is a rather forbidding exercise.

The third approach introduces a^x as solution of the initial value problem

$$\begin{aligned} f'(x) &= kf(x), \\ f(x_0) &= y_0. \end{aligned}$$

We first consider the case $k = 1, x_0 = 0$.

Let $f(x) = A_0 + A_1x + A_2x^2 + A_3x^3 + \dots + A_nx^n$ which we differentiate to get $f'(x) = A_1 + 2A_2x + 3A_3x^2 + 4A_4x^3 + \dots + nA_nx^{n-1}$. Substituting in the differential equation, neglecting the term A_nx^n in $f(x)$ since it is small when x is near 0 and identifying the coefficients we obtain:

$$\begin{aligned} A_0 &= y_0 \quad (\text{from the initial condition}) \\ A_1 &= A_0 \\ 2A_2 &= A_1 \\ 3A_3 &= A_2 \\ &\dots \\ nA_n &= A_{n-1} \end{aligned}$$

from which we get $A_n = \frac{y_0}{n!}$ and $f(x) = y_0 \sum_{i=0}^{i=n} \frac{x^i}{i!} + \dots$ that is $f(x) = f(x_0) \sum_{i=0}^{i=n} \frac{x^i}{i!} + \dots$

It is interesting to note that many of the properties of the exponential function can be recovered from this approximation. In particular, we have an addition formula for a and b near 0 :

$$\begin{aligned} f(a)f(b) &= y_0 \left[1 + a + \frac{a^2}{2!} + \frac{a^3}{3!} + \dots \right] y_0 \left[1 + b + \frac{b^2}{2!} + \frac{b^3}{3!} + \dots \right] \\ &= y_0^2 \left[1 + a + \frac{a^2}{2!} + \frac{a^3}{3!} + \dots \right. \\ &\quad \left. + b + ab + \frac{a^2b}{2!} + \dots \right. \\ &\quad \left. + \frac{b^2}{2!} + \frac{ab^2}{2!} + \dots \right. \\ &\quad \left. + \frac{b^3}{3!} + \dots \right] \\ &= y_0^2 \left[1 + (a+b) + \frac{(a+b)^2}{2!} + \frac{(a+b)^3}{3!} + \dots \right] \\ &= y_0 f(a+b) \end{aligned}$$

since '...' stands for finite remainders and *not* for infinite tails. We thus have

$$\begin{aligned} f(x_0 + h) &= 1 + (x_0 + h) + \frac{(x_0 + h)^2}{2!} + \frac{(x_0 + h)^3}{3!} + \dots \\ &= \left[1 + x_0 + \frac{x_0^2}{2!} + \frac{x_0^3}{3!} + \dots \right] \left[1 + h + \frac{h^2}{2!} + \frac{h^3}{3!} + \dots \right] \\ &= f(x_0)f(h) \end{aligned}$$

which we use to localize.

We can find an approximate solution near x_0 by localizing. Writing $f(x_0+h) = f_{x_0}(h)$ and since by the chain rule $f'(x) = f'_{x_0}(h)$, the differential problem becomes

$$\begin{aligned} f'_{x_0}(h) &= f'_{x_0}(h) \\ f_{x_0}(0) &= y_0 \end{aligned}$$

which gives $f_{x_0}(h) = y_0 \sum_{i=0}^{i=n} \frac{h^i}{i!} + \dots$

Once the existence of a solution $f(x)$ of the initial value problem $f'(x) = f(x)$, $f(0) = 1$ is assumed, the properties of $f(x)$ are easily obtained.

1. $f(x) \neq 0$ for all x . Differentiating $f(x)f(-x)$ we get 0 so that $f(x)f(-x) = c$ with $c = 1$ by the initial condition.

2. $f(-x) = f(x)^{-1}$.

3. Uniqueness. Let g be another solution. Then $[g/f]' = 0$ so that $g = kf$ for some k . From the initial condition, $k = 1$ and $g = f$.

4. Positivity. From **1.** and the differential equation, $f'(x)$ cannot have a zero so that, by the Intermediate Value Theorem, $f'(x)$ must keep the same sign for all x and since $f'(0) = 1$, $f'(x) > 0$ for all x .

5. Increasingness. Follows from **4.**

6. $f(a+b) = f(a)f(b)$. Consider the function $g(a+x)$. Then $g'(a+x) = f'(a+x) = f(a+x) = g(a+x)$, so that $g(x) = kf(x)$ with k such that $k = g(0) = f(a)$. so, $f(a+x) = f(a)f(x)$ for all x .

7. $f(x) = e^x$. We have $f(na) = f(a^n)$ for all positive integer n because it is true for $n = 1$ and, assuming it for n , we have $f((n+1)a) = f(na+a) = f(na)f(a) = f^n(a)f(a) = f^{n+1}(a)$. Defining $e = f(1)$ gives $f(n) = e^n$. Since f is strictly increasing, we have $1 < e$ from $f(0) < f(1)$. We also have from **2.** that $f(-n) = f(n)^{-1} = e^{-n}$ and the result follows.

8. Graph. Since $e > 1$, write $e = 1 + b$ with $b > 0$ so that $e^n = (1 + b)^n \geq 1 + nb$. Since e^x is strictly increasing, $e^x \hat{a} \hat{e}$ when $x \hat{a} \hat{e}$. Finally, $e^{-x} \hat{a} 0$ when $x \hat{a} \hat{e}$ so that $e^x \hat{a} 0$ when $x \hat{a} -\hat{e}$. This gives the qualitative look of the graph.

9. Comparison with power functions: $\lim_{n \hat{a} \hat{e}} (x^n/e^x) = 0$. First show that $\lim_{n \hat{a} \hat{e}} (n/c^n) = 0$, $c > 1$, by setting $c = 1 + b$ and observing that $(1 + b)^n \geq 1 + nb + n(n-1)/2 \hat{u} b^2$ and dividing by n . Now let $\varphi(x) = x/e^n$. Then $\varphi'(x) = e^x(1-x)/e^{2x}$ and $\varphi'(x) < 0$ when $x > 1$. Hence φ is strictly decreasing and $\varphi(x) \hat{a} 0$ when $x \hat{a} \hat{e}$. A similar proof gives $\lim_{n \hat{a} \hat{e}} (x^n/e^x) = 0$.

The logarithm function can then be introduced as inverse of the exponential function. See [2] for details.

Note, however, that $\ln(x)$ is easily approximated near $x_0 = 1$ as solution of $f'(x) = 1/x$ with the initial condition $f(1) = 0$ since, by the chain rule, $f'(x)|_{x_0+h} = f'_h(x_0+h)$ so that if we set $f(x_0+h) = A_0 + A_1h + A_2h^2 + A_3h^3 + \dots$ we have $f'(x)|_{x_0+h} = A_1 + 2A_2h + 3A_3h^2 + \dots$ which must then be equal to $1/[1+h] = 1 - h + h^2 - h^3 + \dots$, that is $A_1 + 2A_2h + 3A_3h^2 + \dots = 1 - h + h^2 - h^3 + \dots$ from which we get the coefficients A_0, A_1, A_2, \dots . In fact, here again, we need not expand near 1 but we can expand it near any $x_0 \neq 0$ as soon as we have $f(x_0)$. The process can therefore be iterated.

13. TRIGONOMETRIC FUNCTIONS. We can define sine and cosine as solutions of the system $f' = g$ and $g' = -f$ with the initial conditions $f(0) = 0$ and $g(0) = 1$ and we can again recover the usual properties from the system. See [2] for the details.

In a more "physical" manner, we can also define them as solutions of $f'' = -f$ with the appropriate initial conditions: $f(0) = 1$ and $f'(0) = 0$ for the cosine and $f(0) = 0$ and $f'(0) = 1$ for the sine and we can again recover the usual properties from the system. In either case, $\pi/2$ is defined as the smallest zero of $\cos x$. See [1] for the details.

[1] R. L. Finney – D. R. Ostbey. *Elementary Differential Equations with Linear Algebra*. Addison Wesley, Reading, 1984.

[2] S. Lang, *Analysis I*. Addison-Wesley, Reading, 1976.

[3] H. Levi, *Polynomials, Power Series and Calculus*. Van Nostrand, Princeton, 1968.